

Overview

Problem: How to obtain an efficient procedure to compute an approximation of the Gromov Wasserstein cost?

Contributions:

- We show first that a **low-rank factorization** (or approximation) of the two input **cost matrices** that define GW, one for each measure, can be exploited to lower the complexity of the **entropic GW** problem from **cubic to quadratic**.
- We show next, independently, that by imposing a **low-nonnegative rank** on the **couplings** involved in the GW problem we obtain a solver only requiring $\mathcal{O}(n^2)$ operations with no prior assumption on input cost matrices.
- Finally, we show that both **low-rank assumptions** (on costs and couplings) can be combined to shave yet another factor and reach a **linear $\mathcal{O}(n)$ GW approximation**.
- We show **experimentally** the efficiency of our approach.

Entropic Gromov Wasserstein

Discrete Distributions: $\mu = \sum_{i=1}^n a_i \delta_{x_i}$ and $\nu = \sum_{j=1}^m b_j \delta_{y_j}$

Cost matrices: $A = [d_{\mathcal{X}}(x_i, x_j)]_{1 \leq i, j \leq n}$, $B = [d_{\mathcal{Y}}(y_i, y_j)]_{1 \leq i, j \leq m}$

Definition of Entropic Gromov Wasserstein:

$$GW_{\varepsilon}((a, A), (b, B)) := \min_{P \in \mathbb{R}_+^{n \times m}} \mathcal{E}_{A, B}(P) - \varepsilon H(P)$$

$$P \mathbf{1}_m = a, P^T \mathbf{1}_n = b$$

where $\mathcal{E}_{A, B}(P) := \sum_{i, j, i', j'} |A_{i, i'} - B_{j, j'}|^2 P_{i, j} P_{i', j'}$

Shannon entropy

Mirror Descent Scheme:

Init: $a, A, b, B, \varepsilon, P$

for $t = 0, \dots, T$:

$C = -4APB$ ← Update the cost: $\mathcal{O}(nm(n+m))$

$K = \exp(-C/\varepsilon)$ ← Update the kernel: $\mathcal{O}(nm)$

$P = \arg \min_{P \geq 0, P \mathbf{1}_m = a, P^T \mathbf{1}_n = b} \text{KL}(P, K)$ ← Solve the entropic OT: $\mathcal{O}(nm)$

Low-rank Costs

Idea: Replace A by $\tilde{A} = A_1 A_2^T$ where $(A_1, A_2) \in (\mathbb{R}_+^{n \times d}) \times (\mathbb{R}_+^{n \times d'})$

Replace B by $\tilde{B} = B_1 B_2^T$ where $(B_1, B_2) \in (\mathbb{R}_+^{n \times d'}) \times (\mathbb{R}_+^{n \times d'})$

→ Updating the cost $C = -4A_1 A_2^T P B_1 B_2^T$ requires now $nm(d+d') + dd'(n+m)$ operations.

Examples:

- SE distance: $A = \left[\|x_i - x_j\|_2^2 \right]_{i, j} = A_1 A_2^T$ with $z = (X^{\odot 2})^T \mathbf{1}_d$,
 $A_1 = [z, \mathbf{1}_n, -2X^T] \in \mathbb{R}^{n \times (d+2)}$, $A_2 = [\mathbf{1}_n, z, X^T] \in \mathbb{R}^{n \times (d+2)}$
 computable in $\mathcal{O}(n)$

- General Distance Matrix: $\|A - A_1 A_2^T\|_F^2 \leq \|A - C_d\|_F^2 + \gamma \|A\|_F^2$

Low-Rank Couplings

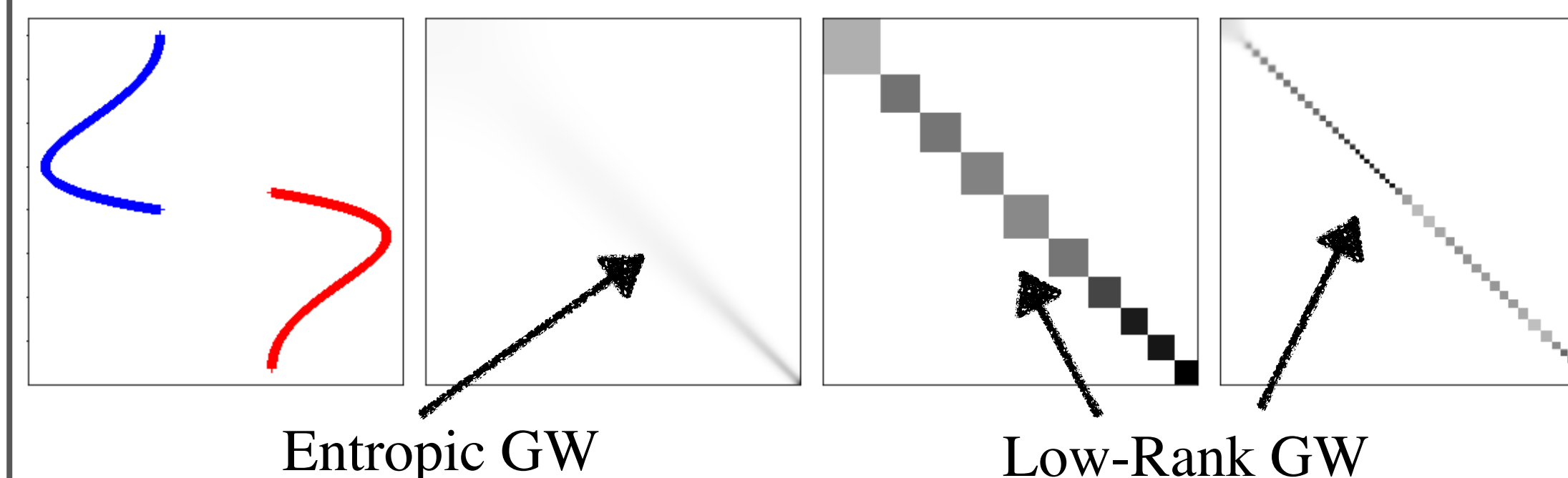
NN-rank: $\text{rk}_+(M) := \min \left\{ q \mid M = \sum_{i=1}^q R_i, \forall i, \text{rk}(R_i) = 1, R_i \geq 0 \right\}$

Low-NN rank couplings:

$$\Pi_{a, b}(r) := \{P \in \mathbb{R}_+^{n \times m} \text{ s.t. } P \mathbf{1}_m = a, P^T \mathbf{1}_n = b \text{ and } \text{rk}_+(P) \leq r\}$$

Definition of Low-Rank Gromov Wasserstein:

$$GW\text{-LR}_r((a, A), (b, B)) := \min_{P \in \Pi_{a, b}(r)} \mathcal{E}_{A, B}(P)$$



Reparametrization of GW-LR:

$$GW\text{-LR}_r((a, A), (b, B)) = \min_{(Q, R, g) \in \mathcal{E}_1(a, b, r) \cap \mathcal{E}_2(r)} \mathcal{E}_{A, B}(Q \text{Diag}(1/g) R^T)$$

$$\mathcal{E}_1(a, b, r) := \{(Q, R, g) \in \mathbb{R}_+^{n \times r} \times \mathbb{R}_+^{m \times r} \times (\mathbb{R}_+^r) \text{ s.t. } Q \mathbf{1}_r = a, R \mathbf{1}_r = b\}$$

$$\mathcal{E}_2(r) := \{(Q, R, g) \in \mathbb{R}_+^{n \times r} \times \mathbb{R}_+^{m \times r} \times (\mathbb{R}_+)^r \text{ s.t. } Q^T \mathbf{1}_n = R^T \mathbf{1}_m = g\}$$

Mirror-Descent Scheme

Algorithm 2: Low-Rank GW

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1 Inputs:  $a, A, B, b, r, Q, R, g$ 
2 for  $t = 1, \dots$  do
3    $C_1 \leftarrow -AQ \text{diag}(1/g)$ 
4    $C_2 \leftarrow R^T B$ 
5    $K^{(1)} \leftarrow Q \odot \exp(4\gamma C_1 C_2 R \text{diag}(1/g))$ 
6    $K^{(2)} \leftarrow R \odot \exp(4\gamma (C_1 C_2)^T Q \text{diag}(1/g))$ 
7    $\omega \leftarrow \mathcal{D}(Q^T C_1 C_2 R)$ 
8    $K^{(3)} \leftarrow g \odot \exp(-4\gamma \omega / g^2)$ 
9    $Q, R, g \leftarrow \underset{\zeta \in \mathcal{C}_1(a, b, r) \cap \mathcal{C}_2(r)}{\text{argmin}} \text{KL}(\zeta, \mathbf{K})$ 
10 end
11 Return:  $\mathcal{E}$ 

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Update the costs: $(n^2 + m^2)r$

Update the kernels: $\mathcal{O}((n+m)r^2)$

Solve the convex Barycenter problem with Dykstra: $\mathcal{O}((n+m)r)$

Double Low-Rank GW

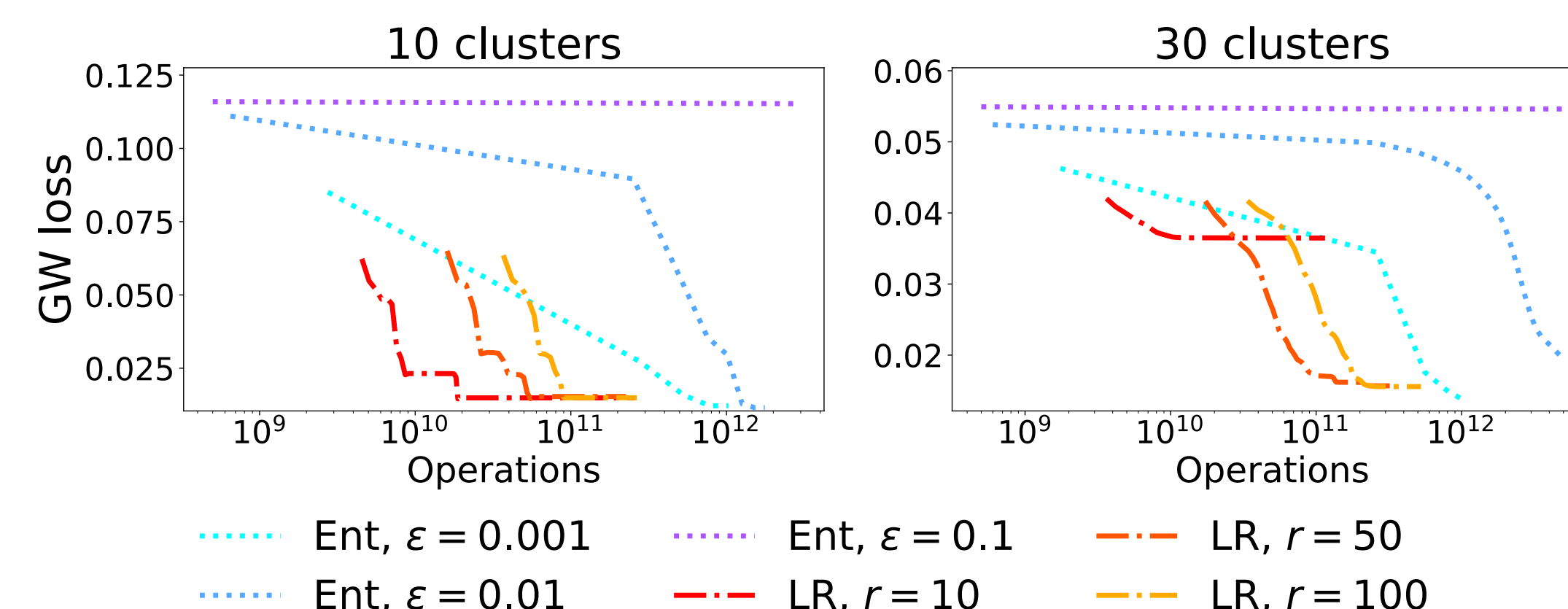
▲ The only steps which remain quadratic are the updates of the costs C_1 and C_2 .

→ By replacing A by $\tilde{A} = A_1 A_2^T$ and B by $\tilde{B} = B_1 B_2^T$:

$$C_1 = -A_1 A_2^T Q \text{Diag}(1/g) \quad \text{and} \quad C_2 = -R^T B_2 B_1^T$$

$$\mathcal{O}(nrd + mrd')$$

Experiments



Comparison of the time-accuracy tradeoff between our method and the Entropic GW. We plot $n = 5000$ samples from two isotropic Gaussian Blobs in 10 and 15-D. We observe that our method obtains **similar GW loss**, while being orders of magnitude **faster**.